

# **STL - Principles and Practice**

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# Agenda

#### Part 0: STL Intro.



#### **Part 1: Containers and Iterators**



#### Part 2: STL Function Objects and Utilities



Part 3-4: STL Algorithms Principles and Practice



# Part 3: STL Algorithms - Principles and Practice

"Prefer algorithm calls to hand-written loops." - Scott Meyers, "Effective STL"

## Correctness

Fewer opportunities to write bugs (less code => less bugs) like:

- iterator invalidation
- copy/paste bugs
- iterator range bugs
- loop continuations or early loop breaks
- guaranteeing loop invariants
- issues with algorithm logic

Code is a liability: maintenance, people, knowledge, dependencies, sharing, etc.

**More code** => more bugs, more test units, more maintenance, more documentation

# Code Clarity

- Algorithm **names** say what they do.
- Raw "for" loops don't (without reading/understanding the whole body).
- We get to program at a higher level of **abstraction** by using well-known **verbs** (find, sort, remove, count, transform).
- A piece of code is **read** many more times than it's **modified**.
- **Maintenance** of a piece of code is greatly helped if all future programmers understand (with confidence) what that code does.

## Modern C++ (C++11/14 standards)

- Modern C++ adds more useful algorithms to the STL library.
- Makes existing algorithms much easier to use due to simplified language syntax and lambda functions (closures).

```
for(vector<string>::iterator it = v.begin(); it != v.end(); ++it) { ... }
```

```
for(auto it = v.begin(); it != v.end(); ++it) { ... }
```

```
for(auto it = v.begin(), end = v.end(); it != end; ++it) { ... }
```

```
std::for_each(v.begin(), v.end(), [](const auto & val) { ... });
```

```
for(const auto & val : v) { ... }
```

# Performance / Efficiency

- Vendor implementations are highly **tuned** (most of the times).
- Avoid some unnecessary temporary copies (leverage **move** operations for objects).
- Function helpers and functors are **inlined** away (no abstraction penalty).
- Compiler optimizers can do a better job without worrying about **pointer aliasing** (autovectorization, auto-parallelization, loop unrolling, dependency checking, etc.).

# The difference between Efficiency and Performance

Why do we care ?

Because: "Software is getting slower more rapidly than hardware becomes faster."

"A Plea for Lean Software" - Niklaus Wirth

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NIKLAUS WIRTH

Efficiency	Performance	let test cad
the amount of work you need to do	how fast you can do that work	strat rational langes
governed by your algorithm	governed by your data structures	

Efficiency and performance are **not dependant** on one another.

# Optimization

Optimization strategy:

- 1. **Identification**: profile the application and identify the worst performing parts.
- 2. **Comprehension**: understand what the code is trying to achieve and why it is slow.
- 3. **Iteration**: change the code based on step 2 and then re-profile; repeat until fast enough.

Very often, code becomes a bottleneck for one of four reasons:

- It's being called too often.
- It's a bad choice of algorithm:  $O(n^2)$  vs O(n), for example.
- It's doing unnecessary work or it is doing necessary work too frequently.
- The data is bad: either too much data or the layout and access patterns are bad.



# **Generic Programming Drawbacks**

- abstraction penalty
- implementation in the interface
- early binding
- horrible error messages (*no formal specification* of interfaces, **yet**)
- duck typing
- algorithm could work on some data types, but fail to work/compile on some other new data structures (different iterator category, no copy semantics, etc)

We need to fully specify requirements on algorithm types => **Concepts** 

# What Is A Concept, Anyway ?

*Formal* specification of concepts makes it possible to **verify** that *template arguments* satisfy the **expectations** of a template or function during overload resolution and template specialization.

Examples from **STL**:

- DefaultConstructible, MoveConstructible, CopyConstructible
- MoveAssignable, CopyAssignable,
- Destructible
- EqualityComparable, LessThanComparable
- Predicate, BinaryPredicate
- Compare
- FunctionObject
- Container, SequenceContainer, ContiguousContainer, AssociativeContainer
- Iterator
  - InputIterator, OutputIterator
  - ForwardIterator, BidirectionalIterator, RandomAccessIterator

### **Template Constraints Using C++17 Concepts**

#### An example: Balanced reduction

```
template<ForwardIterator I, BinaryOperation Op>
  requires EqualityComparable<ValueType<I>, Domain<Op>>()
Domain<Op> reduce(I it, DistanceType<I> n, Op op)
// precondition: n != 0, "op" is associative
  if (n == 1)
    return *it;
  DistanceType<I> h = n / 2;
  return op ( reduce (it, h, op),
             reduce(it + h, n - h, op));
```

\*\*\* For a better/efficient implementation of a generic **reduce**, see the longer (complex) implementation from *Elements of Programming*, by Alexander Stepanov.

## Compare Concept

Why is this one special ? Because ~50 STL facilities (algorithms & data structures) expect a *Compare* type.

template< class RandomIt, class Compare >
void sort( RandomIt first, RandomIt last, Compare comp );

Concept relations:

Compare << BinaryPredicate << Predicate << FunctionObject << Callable

A type satisfies Compare if:

- it satisfies *BinaryPredicate* bool comp(\*iter1, \*iter2);
- it establishes a *strict weak ordering* relationship

Irreflexivity	$\forall$ a, comp(a,a) == false
Antisymmetry	<pre>∀ a, b, if comp(a,b) == true =&gt; comp(b,a) == false</pre>
Transitivity	<pre>∀ a, b, c, if comp(a,b) == true and comp(b,c) == true =&gt; comp(a,c) == true</pre>

{ partial ordering }

```
vector<string> v = { ... };
sort(v.begin(), v.end());
sort(v.begin(), v.end(), less<>());
sort(v.begin(), v.end(), [](const string & s1, const string & s2)
ł
  return s1 < s2;
});
sort(v.begin(), v.end(), [](const string & s1, const string & s2)
{
  return stricmp(s1.c str(), s2.c str()) < 0;</pre>
});
```

```
struct Point { int x; int y; };
vector<Point> v = { ... };
sort(v.begin(), v.end(), [](const Point & p1, const Point & p2)
{
  return (p1.x < p2.x) && (p1.y < p2.y);
});
```

Is this a good Compare predicate for 2D points ?

```
Definition:
if comp(a,b) == false && comp(b,a) == false
=> a and b are equivalent
```

```
Let { P1, P2, P3 }
x1 < x2; y1 > y2;
x1 < x3; y1 > y3;
x2 < x3; y2 < y3;
```

#### =>

P1 11 9 33 72 X3 0 Y2 X1

P2 and P1 are unordered (P2 ?P1) comp(P2,P1) == false && comp(P1,P2) == false P1 and P3 are unordered (P1 ?P3) comp(P1,P3) == false && comp(P3,P1) == false P2 and P3 are ordered (P2 <P3) comp(P2,P3) == true && comp(P3,P2) == false

#### =>

P2 is **equivalent** to P1 P1 is **equivalent** to P3 P2 is **less than** P3

### Compare Concept

**Partial ordering** relationship: Irreflexivity + Antisymmetry + Transitivity

*Strict weak ordering* relationship: *Partial ordering* + *Transitivity of Equivalence* 

*Total ordering* relationship: *Strict weak ordering* + equivalence must be the same as equality

Irreflexivity	$\forall$ a, comp(a,a) == false
Antisymmetry	<pre>∀ a, b, if comp(a,b) == true =&gt; comp(b,a) == false</pre>
Transitivity	<pre>∀ a, b, c, if comp(a,b)=true and comp(b,c)==true =&gt; comp(a,c)==true</pre>
Transitivity of equivalence	<b>if a</b> is equivalent to <b>b</b> and <b>b</b> is equivalent to <b>c</b> => <b>a</b> is equivalent to <b>c</b>

```
struct Point { int x; int y; };
vector<Point> v = { ... };
sort(v.begin(), v.end(), [](const Point & p1, const Point & p2)
{
   return (p1.x * p1.x + p1.y * p1.y) <
        (p2.x * p2.x + p2.y * p2.y);
});
```

Is this a good Compare predicate for 2D points ?

```
struct Point { int x; int y; };
vector<Point> v = { ... };
sort(v.begin(), v.end(), [](const Point & p1, const Point & p2)
{
    if (p1.x < p2.x) return true;
    if (p2.x < p1.x) return false;
    return p1.y < p2.y;
});
```

Is this a good Compare predicate for 2D points ?

The general idea is to pick an **order** in which to compare **elements/parts** of the object. (in our example we first compared by **x** coordinate, and then by **y** coordinate for equivalent **x**)

This strategy is analogous to how a **dictionary** works, so it is often called *"dictionary order"*, or *"lexicographical order"*.

The STL implements dictionary ordering in at least three places:

**std::pair<T, U>** - defines the six comparison operators in terms of the corresponding operators of the pair's components

std::tuple< ... Types> - generalization of pair

std::lexicographical\_compare() algorithm

- Two ranges are compared element by element
- The first mismatching element defines which range is lexicographically *less* or *greater* than the other

• ...

# **Prefer Member Functions To Similarly Named Algorithms**

The following member functions are available for *associative containers*:

- .count()
- .find()
- .equal\_range()
- .lower\_bound() // only for ordered containers
- .upper\_bound() // only for ordered containers

The following member functions are available for *list containers*:

- .remove() .remove\_if()
- .unique()
- .sort()
- .merge()
- .reverse()

These member functions are always **faster** than their similarly named generic algorithms.

Why? They can leverage the *implementation details* of the underlying data structure.

#### **Prefer Member Functions To Similarly Named Algorithms**

```
set<string> s = {...}; // 1 million elements
```

```
// worst case: 40 comparisons, average: 20 comparisons
auto it = s.find("stl");
if (it != s.end()) {...}
```

```
// worst case: 1 million comparisons, average: ½ million comparisons
auto it = std::find(s.begin(), s.end(), "stl");
if (it != s.end()) {...}
```

Why?

### **Prefer Member Functions To Similarly Named Algorithms**

#### std::list<> specific algorithms

std::sort() doesn't work on lists (Why?)
=> call .sort() member function

.remove() and .remove\_if() don't need to use the erase/remove idiom. They directly remove matching elements from the list.

.remove() and .remove\_if() are more efficient than the generic algorithms, because they just relink nodes with the need to copy or move elements.

#### **Binary search operations (on sorted ranges)**

binary\_search() // helper (incomplete interface - Why ?)
lower\_bound() // returns an iter to the first element not less than the given value
upper\_bound() // returns an iter to the first element greater than the certain value

equal\_range() = { lower\_bound(), upper\_bound() }

```
// properly checking return value
auto it = lower_bound(v.begin(), v.end(), 5);
if ( it != v.end() && (*it == 5) )
{
    // found item, do something with it
}
else // not found, insert item at the correct position
{
    v.insert(it, 5);
}
```

#### **Binary search operations (on sorted ranges)**

#### Counting elements equal to a given value

```
vector<string> v = { ... }; // sorted collection
size_t num_items = std::count(v.begin(), v.end(), "stl");
```

Instead of using std::count() generic algorithm, use binary search instead.

```
auto range = std::equal_range(v.begin(), v.end(), "stl");
size t num items = std::distance(range.first, range.second);
```

# **Extend STL With Your Generic Algorithms**

```
Eg.
```

```
template<class Container, class Value>
void name_this_algorithm(Container & c, const Value & v)
{
    if ( find(begin(c), end(c), v) == end(c) )
        c.emplace_back(v);
        assert( !c.empty() );
}
```

## **Extend STL With Your Generic Algorithms**

```
Eg.
```

```
template<class Container, class Value>
bool erase if exists (Container & c,
                     const Value & v)
  auto found = std::find(begin(c), end(c), v);
  if (found != end(v))
    c.erase(found); // call 'erase' from STL container
    return true;
  return false;
```

#### **Consider Adding Range-based Versions of STL Algorithms**

```
template< class InputRange, class T > inline
typename auto find(InputRange && range, const T & value)
 return std::find(begin(range), end(range), value);
template< class InputRange, class UnaryPredicate > inline
typename auto find if(InputRange && range, UnaryPredicate pred)
 return std::find if (begin (range), end (range), pred);
template< class RandomAccessRange, class BinaryPredicate > inline
void sort(RandomAccessRange && range, BinaryPredicate comp)
  std::sort(begin(range), end(range), comp);
```

namespace **range** {

#### **Consider Adding Range-based Versions of STL Algorithms**

```
vector<string> v = \{ \dots \};
auto it = range::find(v, "stl");
string str = *it;
auto chIt = range::find(str, 't');
auto it2 = range::find if(v, [](const auto & val) { return val.size() > 5; });
range::sort(v);
range::sort(v, [](const auto & val1, const auto & val2)
                { return val1.size() < val2.size(); } );</pre>
```

# **STL Abuse**

Please don't code like this !

Extract algorithm **intermediate** results into **named** local variables (iterators, values, etc.)