Open4Tech Summer School 2021

So You Think You Can # Building an app using Blockchain A list<> of data structures you should add in your learning queue<> Build a video call react app with WebRTC and Socket.io Importance of good coding habits

Build cryptocurrencies exchange using React Why we code Getting a11y right State of the Art Natural Language Processing for Noobs Java Swing Crash Course

28 iunie - 16 iulie 2021 http://inf.ucv.ro/~ summer-school/



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	Luni	Marti	Miercuri	Joi	Vineri
	28 iunie	29 iunie	30 iunie	1 iulie	2 iulie
2-4pm	So You Think You Can # (hashing algorithms & containers)	So You Think You Can # (hashing algorithms & containers)	A list<> of data structures you should add in your learning queue<>	A list<> of data structures you should add in your learning queue<>	Building an app using Blockchain
4-6pm	Why we code	Why we code	Getting ally right	Java Swing Crash Course	
	5 iulie	6 iulie	7 iulie	8 iulie	9 iulie
2-4pm	Build a video call react app with WebRTC and Socket.io	Build a video call react app with WebRTC and Socket.io	Build a video call react app with WebRTC and Socket.io	State of the Art Natural Language Processing for Noobs	State of the Art Natural Language Processing for Noobs
4-6pm					
	12 iulie	13 iulie	14 iulie	15 iulie	16 iulie
2-4pm	Build cryptocurrencies exchange using React	Build cryptocurrencies exchange using React	Code conventions: Importance of good coding habits	Code conventions: Importance of good coding habits	
4-6pm					



So You Think You Can

Hashing Algorithms and Containers

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Victor Ciura Principal Engineer



June 2021





So You Think You Can

Not a C# workshop

Abstract

Most programming languages offer some kind of associative "arrays" or containers. They may be called differently: maps, dictionaries, hash-maps, unordered-maps, hash-tables, etc.

If you've never heard of them, this workshop is for you. If you've seen them before, but you're not sure which one to use for a particular task, this workshop is for you. If you're confident in using such fast lookup structures, great! But you'll still be surprised by the details we're going to cover in this lecture.

After this journey from the very basics of lookup data structures up to advanced hashing techniques, you'll feel more confident when & how to use them effectively.

This workshop assumes familiarity with the C++ language, as we'll focus on hashing facilities in the standard library and beyond. We'll take a deep dive into hashing algorithms and hashed data structures, both in design and examples.











What they teach you

initial call to sort entire array Quicksort(A, 1, length[A])

```
Quicksort (A, p, r)
if p < r
then q = Partition(A, p, r)
Quicksort(A, p, q)
Quicksort(A, q+1, r)</pre>
```

```
Partition (A, p, r)

x = A[p]

i = p-1

j = r+1

while TRUE

do repeat j = j-1

until A[j] <= x

repeat i = i+1

until A[i] >= x

if i < j

then exchange A[i], A[j]

else return j
```



O(n log n)

The Big-O

Algorithm	Data structure	Time complexity:Best	Time complexity:Average	Time complexity:Worst	Space complexity:Worst
Quick sort	Array	O(<i>n</i> log(<i>n</i>))	O(<i>n</i> log(<i>n</i>))	O(<i>n</i> ²)	O(<i>n</i>)
Merge sort	Array	O(<i>n</i> log(<i>n</i>))	$O(n \log(n))$ $O(n \log(n))$ $O(n \log(n))$ $O(n \log(n))$		O(<i>n</i>)
Heap sort	Array	O(<i>n</i> log(<i>n</i>))	O(<i>n</i> log(<i>n</i>))	O(<i>n</i> log(<i>n</i>))	O(1)
Smooth sort	Array	O(<i>n</i>)	O(<i>n</i> log(<i>n</i>))	O(<i>n</i> log(<i>n</i>))	O(1)
Bubble sort	Array	O(<i>n</i>)	O(<i>n</i> ²)	O(<i>n</i> ²)	O(1)
Insertion sort	Array	O(<i>n</i>)	O(<i>n</i> ²)	O(<i>n</i> ²)	O(1)
Selection sort	Array	O(<i>n</i> ²)	O(<i>n</i> ²)	O(<i>n</i> ²)	O(1)
Bogo sort	Array	O(<i>n</i>)	O(<i>n n</i> !)	O(∞)	O(1)

wikipedia.org/wiki/Computational_complexity_theory

What about Data Structures ? 爹



Data structures along with the operations they provide, also have complexity guarantees



STL Containers Big-O cheat-sheet

	A	В	C	D	E	F	G	н	1
1	C++ STL	insert @end	insert @pos	erase @end	erase @pos	find	sort	iterator	comment
2	vector	0(1)	O(dist(pos,end))	0(1)	O(dist(pos,end))	0(n)	O(n*log(n))	RandomAccess	array
3	dequeue	<pre>@begin/@end 0(1)</pre>	O(dist(pos,begin/end))	<pre>@begin/@end 0(1)</pre>	<pre>O(dist(pos,begin/end))</pre>	0(n)	O(n*log(n))	RandomAccess	
4	list	0(1)	0(1)	0(1)	@pos 0(1); @key 0(n)	0(n)	O(n*log(n))	Bidirectional	doubly linked
5	stack	0(1) push()	-	O(1) pop()	-	0(n)		same as container	adaptor <dequeue, list,="" vector=""></dequeue,>
6	queue	0(1) push()	-	O(1) pop() @begin	27	0(n)	0712	same as container	adaptor <dequeue, list=""></dequeue,>
7	set/map	-	O(log(n))	100	<pre>@pos 0(1); @key 0(log(n)+count(key))</pre>	0(log(n))	sorted	Bidirectional	red-black tree (balanced BST)
8	unordered_set/ unordered_map	-	avg O(1); worst O(n)	3 -	<pre>@pos avg 0(1) worst 0(n); @key 0(count(key))</pre>	avg O(1); worst O(n)	-	Forward	hash_set/hash_map
9	priority_queue	<pre>push() 0(log(n))</pre>	-	<pre>pop() 0(log(n))</pre>	-	top() 0(1)		RandomAccess	<pre>adaptor<vector, dequeue=""> => constant time extraction of the largest (default) element, at the expense of logarithmic insertion</vector,></pre>
10	make_heap(range)	push_heap() O(2*log(n))	1.00	<pre>pop_heap() 0(2*log(n))</pre>	÷	max is first	0(n*log(n))	RandomAccess	constructs a max heap in the range

The difference between Efficiency and Performance

Why do we care ?

Because: "Software is getting slower more rapidly than hardware becomes faster."

"A Plea for Lean Software" - Niklaus Wirth

	Efficiency	Performance	Void, systemate, and parameters treatment of basis and dynamic data	
	the amount of work you need to do	how fast you can do that work	stractores, sanfag, recursion signation, language structures, suid compiling	NIKLAUS WIRTH
	governed by your algorithm	governed by your data structures		Algorithms +
i Effic	ciency and performance are <i>not neces</i>	PRENT CP. HALL BUILDS IN AUTOWATIO COUPL TATIEN	Structures = Programs	

What about Performance ? 💉

How **fast** can the CPU execute **each step** from the algorithms.

This is mostly determined by the native (CPU) **data types** used and your choice of **data structures**.



That's all great and still relevant, but...

















Cache

Cache Memory



80486 (1989) This is the <u>first</u> CPU of this generation which has some **cache** on the CPU. It is a 8KB **unified** cache which means it is used for *data* and *instructions*.



80586 (1993)

The 586 or Pentium-1 uses a a **split** level 1 cache. 8 KB each for data and instructions.

The cache was split so that the *data* and *instruction* caches could be individually tuned for their specific use.

You still have a small yet very fast 1st cache near the CPU, and a larger but slower 2nd cache on the *motherboard*.



80686 (1995)

The 686 or Pentium Pro chip, depending on the model, had a 256Kb, 512KB or 1MB on board cache.

Half the space in the chip is used by the cache.



Pentium 2 (1997)

For economy reasons the 2nd cache is **not** in the CPU.

CPU package is on a PCB with separate chips for:

- CPU (and 1st cache)
- 2nd cache





 Pentium
 III
 (1999)

 Pentium
 4
 (2000)

As technology progresses and we start put create chips with smaller components it gets financially possible to put the 2nd cache back in the actual CPU die.

However there is still a split:

Very fast 1st cache snuggled up to the CPU. With one 1st cache per CPU core and a larger but less fast 2nd cache next to the core.



L1 and L2 caches not enough...

=>L3 cache

Nehalem 1st gen Core i7 series (2008)



Cache Latency

Core i7 Xeon 5500 Data Source Latency

local	L1 CACHE hit,		~4	cycles	(2.1 -	1.2 ns
local	L2 CACHE hit,		~10	cycles	(5.3 -	3.0 ns
local	L3 CACHE hit,	line unshared	~40	cycles	(21.4 -	12.0 ns
local	L3 CACHE hit,	shared line in another core	~65	cycles	(34.8 -	19.5 ns
local	L3 CACHE hit,	modified in another core	~75	cycles	(40.2 -	22.5 ns

local DRAM remote DRAM ~60 ns ~100 ns

- Memory Layout
- Memory Access Patterns

Container access patterns

Jumping around through memory,

chasing pointers...



Further Study



ciura.ro/presentations/2021/Open4Tech/ChasingNodes.pdf

Data structures... everywhere

SPACE EAST" LIGHTYE.



cppreference.com/w/cpp/container

A List<> Of Data Structures You Should Add In Your Learning Queue<>

In this workshop we'll explain the mechanics behind data structures, and we'll deep dive into the meaning and usage of the most common ones.

Let's try together to get some insights in some data structures and their pitfalls.

We'll discover, by lots of examples, the strengths and weaknesses of each data structure and find good use cases for:

- Vector and List
- Stack and Queue
- Hash Table
- Graphs and Trees
- Heap.

Difficulty: intermediate

Format: 2 days x 2h



Trainer: Nicolae Telechi

Caphyon, Senior Software Developer

	Luni	Marti	Miercuri	Joi	Vineri
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www.caphyon.ro/open4tech-2021-a-list-of-data-structures-you-should-add-in-your-learning-queue/

90% of situations, a great choice*



* when performance matters

Today's focus

Most programming languages offer some kind of associative

"arrays" or containers.

They may be called differently:

- maps
- dictionaries
- hash-maps
- unordered-maps
- hash-tables

Today's focus



unordered_map

unordered_multiset

unordered_multimap

Hash Functions & Hash Tables

A hash function is any function that can be used to map data of *arbitrary size* to data of *fixed size* (hash code).

Hash functions are used in hash tables, to <u>quickly</u> locate a data record given its search key.
Hash Functions & Hash Tables

The hash function is used to map the search key to an **index;** the index gives the place in the hash table where the corresponding record should be stored/found.

The **domain** of a hash function (the set of possible keys) is larger than its **range** (the number of different table indices), and so it will map several <u>different keys</u> to the <u>same index</u>.

Hash Functions & Hash Tables

Each slot (**bucket**) of a hash table is associated with a <u>set of records</u>, rather than a single record.

Visualize it



Determinism

A hash procedure must be deterministic – meaning that for a given input value it must always generate the <u>same</u> <u>hash</u> value.

Uniformity

A good hash function should map the expected inputs as evenly as possible over its output range.

That is, every hash value in the output range should be generated with roughly the <u>same probability</u>.



It is often desirable that the output of a hash function have <u>fixed size</u>.

If, for example, the output is constrained to 32-bit integer values, the hash values can be used to index into an array (eg. hash tables).

Non-invertible

In *cryptographic* applications, hash functions are typically expected to be practically non-invertible, meaning that it is not realistic to <u>reconstruct</u> the input datum from its hash value alone, without spending great amounts of computing time.

Questions

- How should one <u>combine</u> hash codes from your bases and data members to create a "good" hash function?
- How does one know if you have a good hash function?
- If somehow you knew you had a bad hash function, how would you change it for a type built out of several bases and/or data members?

How does one hash this class?

```
class Customer
{
   std::string firstName;
   std::string lastName;
   int age;
   // ...
};
```

std::hash<Key>

Defined in header <functional>

std::size_t h = std::hash<std::string>{}(firstName);

- Accepts a single parameter of type Key
- Returns a value of type size_t that represents the hash value of the parameter
- Does not throw exceptions when called
- If k1 and k2 are equal => std::hash<Key>()(k1) == std::hash<Key>()(k2)
- If k1 and k2 are different, the *probability* that std::hash<Key>()(k1) == std::hash<Key>()(k2) should be very small, approaching
 1.0/std::numeric_limits<size_t>::max()

std::hash<Key>

Standard specializations for *basic* types:

```
template< class T > struct hash<T*>;
```

template<> struct hash<bool>; template<> struct hash<char>; template<> struct hash<signed char>; template<> struct hash<unsigned char>; template<> struct hash<char16 t>; template<> struct hash<char32 t>; template<> struct hash<wchar t>; template<> struct hash<short>; template<> struct hash<unsigned short>; template<> struct hash<int>; template<> struct hash<unsigned int>; template<> struct hash<long>; template<> struct hash<long long>; template<> struct hash<unsigned long>; template<> struct hash<unsigned long long>; template<> struct hash<float>; template<> struct hash<double>; template<> struct hash<long double>;

Standard specializations for *library* types:

std::hash<std::string>
std::hash<std::wstring>
std::hash<std::unique_ptr>
std::hash<std::shared_ptr>
std::hash<std::bitset>
//...

```
class Customer
  std::string firstName;
  std::string lastName;
  int
              age;
public:
// ...
  std::size t hash code() const
  {
    std::size t k1 = std::hash<std::string>{}(firstName);
    std::size t k2 = std::hash<std::string>{}(lastName);
    std::size t k3 = std::hash<int>{}(age);
    return hash combine(k1, k2, k3); // what algorithm is this?
};
```

Is this a good hash strategy?

What if we wanted to use another hash algorithm?

boost::hash_combine

```
template <class T>
inline void hash_combine(std::size_t & seed, const T & v)
{
    std::hash<T> hasher;
    seed ^= hasher(v) + 0x9e3779b9 + (seed<<6) + (seed>>2);
}
```

The magic number is supposed to be 32 "random" bits:

- each is equally likely to be 0 or 1
- with no simple correlation between the bits

A common way to find a pattern of such bits is to use the binary expansion of an *irrational number*.

In this case, that number is the reciprocal of the golden ratio:

 $\varphi = (1 + \text{sqrt}(5)) / 2$ 2^32 / $\varphi = 0x9e3779b9$

http://stackoverflow.com/questions/35985960/c-why-is-boosthash-combine-the-best-way-to-combine-hash-values

FNV-1A

```
std::size t fnv1a(void const * key, std::size t len)
  std::size t h = 14695981039346656037u;
  unsigned char const * p = static_cast<unsigned char const*>(key);
  unsigned char const * const e = p + len;
  for (; p < e; ++p)
   h = (h ^ *p) * 1099511628211u;
  return h;
```

The FNV hash was designed for fast hash-table and checksum use (not cryptography).

Hash with FNV-1A

```
class Customer
  std::string firstName;
  std::string lastName;
  int
              age;
public:
// ...
  std::size t hash code() const
  {
    std::size t k1 = fnv1a(firstName.data(), firstName.size());
    std::size t k2 = fnv1a(lastName.data(), lastName.size());
    std::size t k3 = fnv1a(&age, sizeof(age));
    return hash combine(k1, k2, k3); // what algorithm is this?
};
```

Ok, but our algorithm is still "polluted" by the combine step...

Anatomy Of A Hash Function

- 1. Initialize internal state.
- 2. Consume bytes into internal state.
- 3. Finalize internal state to result_type (usually size_t).

Anatomy Of A Hash Function

```
std::size_t fnv1a(void const * key, std::size_t len)
{
    std::size t h = 14695981039346656037u; ← initialize internal state
```

```
// consume bytes into internal state:
unsigned char const * p = static_cast<unsigned char const*>(key);
unsigned char const * const e = p + len;
for (; p < e; ++p)
    h = (h ^ *p) * 1099511628211u;
```

return h; ← finalize internal state to size_t

Repackaging this algorithm to make the three stages separately accessible

```
class fnv1a
```

```
std::size_t h = 14695981039346656037u;  ← initialize internal state
public:
```

```
// consume bytes into internal state
 void operator()(void const * key, std::size t len) noexcept
  {
   unsigned char const * p = static cast<unsigned char const*>(key);
   unsigned char const * const e = p + len;
   for (; p < e; ++p)
     h = (h ^ *p) * 1099511628211u;
 explicit operator size t() noexcept \leftarrow finalize internal state to size t
   return h;
};
```

```
class Customer
  std::string firstName;
  std::string lastName;
  int
              age;
public:
 // ...
  std::size t hash code() const
    fnv1a hasher;
    hasher(firstName.data(), firstName.size());
    hasher(lastName.data(), lastName.size());
    hasher(&age, sizeof(age));
    return static cast<std::size t>(hasher); // no more hash combine() !!!
};
```

The same technique can be used with almost every existing hashing algorithm.

```
Now we are using a "pure" FNV-1A algorithm
      for the entire data structure.
```

Combining Types

```
class Sale
 Customer customer;
 Product product;
          date;
 Date
public:
  std::size t hash code() const
    std::size_t h1 = customer.hash code();
    std::size t h2 = product.hash code();
    std::size t h3 = date.hash code();
    return hash combine(h1, h2, h3); // OMG, it's back :(
};
```

How do we use just FNV-1A for the entire class?

Proposal by:

Howard Hinnant, Vinnie Falco, John Bytheway

N3980 / 2014-05-24

```
class Customer
  std::string firstName;
  std::string lastName;
  int
              age;
public:
 // ...
  std::size t hash code() const
    fnv1a hasher;
    hasher(firstName.data(), firstName.size());
    hasher(lastName.data(), lastName.size());
    hasher(&age, sizeof(age));
    return static cast<std::size t>(hasher); // no more hash combine() !!!
};
```

```
class Customer
  std::string firstName;
  std::string lastName;
  int
              age;
public:
 // ...
  friend void hash append(fnv1a & hasher, const Customer & c)
    hasher(c.firstName.data(), c.firstName.size());
    hasher(c.lastName.data(), c.lastName.size());
    hasher(&c.age, sizeof(c.age));
};
```

Let some other piece of code *construct* and *finalize* fnv1a. Customer only *appends* to the state of fnv1a.

```
class Sale
  Customer customer;
  Product product;
           date;
  Date
public:
  friend void hash append(fnv1a & hasher, const Sale & s)
  {
    hash append(hasher, s.customer);
    hash append(hasher, s.product);
    hash append(hasher, s.date);
  }
```

};

Types can recursively build upon one another's hash_append() to build up state in fnv1a object.

```
class Customer
  std::string firstName;
  std::string lastName;
  int
              age;
public:
 // ...
  friend void hash append(fnv1a & hasher, const Customer & c)
  ł
    hash append(hasher, c.firstName);
    hash append(hasher, c.lastName);
    hash append(hasher, c.age);
};
```

Primitive and std-defined types can be given hash_append() overloads => simplified & uniform interface

hash_append() / Abstracting the algorithm

```
class Customer
  std::string firstName;
  std::string lastName;
  int
             age;
public:
 // ...
  template<class HashAlgorithm>
  friend void hash append(HashAlgorithm & hasher, const Customer & c)
  {
    hash append(hasher, c.firstName);
    hash append(hasher, c.lastName);
    hash append(hasher, c.age);
};
```

If all hash algorithms use a *uniform interface*, we can swap any hasher into our data type.

hash_append() / Primitives

For **primitive types** that are **contiguously hashable** we can just send their <u>bytes</u> to the hash algorithm in hash_append().

Eg.

```
template <class HashAlgorithm>
void hash_append(HashAlgorithm & hasher, int i)
{
    hasher(&i, sizeof(i));
}
template <class HashAlgorithm, class T>
void hash_append(HashAlgorithm & hasher, T * p)
{
    hasher(&p, sizeof(p));
}
```

A complicated class is ultimately made up of **scalars** located in discontiguous memory.

hash_append() appends each byte to the HashAlgorithm state by *recursing* down into the data structure to find the scalars.

Prerequisites:

- Every type has a hash_append() overload
- The overload will either call hash_append() on its bases and members, or it will send bytes of its memory to the HashAlgorithm
- No type is aware of the concrete HashAlgorithm type.

How to use hash_append()

HashAlgorithm hasher;

```
hash_append(hasher, my_type);
```

```
return static_cast<size_t>(hasher);
```

Wrap the whole thing up in a conforming hash functor

```
template <class HashAlgorithm>
struct GenericHash
  using result type = typename HashAlgorithm::result type;
 template <class T>
  result type operator()(const T & t) const noexcept
  {
    HashAlgorithm hasher;
    hash append(hasher, t);
    return static cast<result type>(hasher);
};
```

unordered_set<Customer, GenericHash<fnv1a>> my_set;

Change Hashing Algorithms

unordered_set<Sale, GenericHash<fnv1a>> my_set; unordered_set<Sale, GenericHash<SipHash>> my_set; unordered_set<Sale, GenericHash<Spooky>> my_set; unordered_set<Sale, GenericHash<Murmur>> my_set; unordered_set<Sale, GenericHash<CityHash>> my_set;

It becomes trivial to <u>experiment</u> with different hashing algorithms to optimize performance, minimize collisions.

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4-6pm					







So You Think You Can

Hashing Algorithms and Containers

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Victor Ciura Principal Engineer





June 2021



Recap from Part 1



- Determinism
- Uniformity
- Defined range
- Non-invertible
Anatomy Of A Hash Function

- 1. Initialize internal state.
- 2. Consume bytes into internal state.
- 3. Finalize internal state to result_type (usually size_t).

Repackaging this algorithm to make the three stages separately accessible

```
class fnv1a
```

```
std::size_t h = 14695981039346656037u;  ← initialize internal state
public:
```

```
// consume bytes into internal state
 void operator()(void const * key, std::size t len) noexcept
  {
   unsigned char const * p = static cast<unsigned char const*>(key);
   unsigned char const * const e = p + len;
   for (; p < e; ++p)
     h = (h ^ *p) * 1099511628211u;
 explicit operator size t() noexcept \leftarrow finalize internal state to size t
   return h;
};
```

hash_append()

```
class Customer
  std::string firstName;
  std::string lastName;
  int
              age;
public:
 // ...
  friend void hash append(fnv1a & hasher, const Customer & c)
    hasher(c.firstName.data(), c.firstName.size());
    hasher(c.lastName.data(), c.lastName.size());
    hasher(&c.age, sizeof(c.age));
};
```

Let some other piece of code *construct* and *finalize* fnv1a. Customer only *appends* to the state of fnv1a.

hash_append()

```
class Sale
  Customer customer;
  Product product;
           date;
  Date
public:
  friend void hash append(fnv1a & hasher, const Sale & s)
  {
    hash append(hasher, s.customer);
    hash append(hasher, s.product);
    hash append(hasher, s.date);
  }
```

};

Types can recursively build upon one another's hash_append() to build up state in fnv1a object. Wrap the whole thing up in a conforming hash functor

```
template <class HashAlgorithm>
struct GenericHash
  using result type = typename HashAlgorithm::result type;
  template <class T>
  result type operator()(const T & t) const noexcept
  {
    HashAlgorithm hasher;
    hash append(hasher, t);
    return static cast<result type>(hasher);
};
```

std::unordered_set<Customer, GenericHash<fnv1a>> my_set;

Change Hashing Algorithms

unordered_set<Sale, GenericHash<fnv1a>> my_set; unordered_set<Sale, GenericHash<SipHash>> my_set; unordered_set<Sale, GenericHash<Spooky>> my_set; unordered_set<Sale, GenericHash<Murmur>> my_set; unordered_set<Sale, GenericHash<CityHash>> my_set;

It becomes trivial to <u>experiment</u> with different hashing algorithms to optimize performance, minimize collisions.

std::hash<Key>

Defined in header <functional>

std::size_t h = std::hash<std::string>{}(firstName);

- Accepts a single parameter of type Key
- Returns a value of type size_t that represents the hash value of the parameter
- Does not throw exceptions when called
- If k1 and k2 are equal => std::hash<Key>()(k1) == std::hash<Key>()(k2)
- If k1 and k2 are different, the *probability* that std::hash<Key>()(k1) == std::hash<Key>()(k2) should be very small, approaching
 1.0/std::numeric_limits<size_t>::max()

std::hash<Key>

Standard specializations for *basic* types:

template< class T > struct hash<T*>;

template<> struct hash<bool>; template<> struct hash<char>; template<> struct hash<signed char>; template<> struct hash<unsigned char>; template<> struct hash<char16 t>; template<> struct hash<char32 t>; template<> struct hash<wchar t>; template<> struct hash<short>; template<> struct hash<unsigned short>; template<> struct hash<int>; template<> struct hash<unsigned int>; template<> struct hash<long>; template<> struct hash<long long>; template<> struct hash<unsigned long>; template<> struct hash<unsigned long long>; template<> struct hash<float>; template<> struct hash<double>; template<> struct hash<long double>;

Standard specializations for *library* types:

std::hash<std::string>
std::hash<std::wstring>
std::hash<std::unique_ptr>
std::hash<std::shared_ptr>
std::hash<std::bitset>

//...

</Recap>

std::hash<Key>

Standard specializations for *library* types:

```
std::hash<std::string>
std::hash<std::wstring>
```

std::hash<std::unique_ptr>
std::hash<std::shared_ptr>

std::hash<std::bitset>

//...

Exploring string hash tables <code walk-through>

What we want:

"A hash table mapping string keys (case-insensitive) to some custom data type."

Starting point:

```
template<
  class Key,
  class T,
  class Hash = std::hash<Key>,
  class Hash = std::equal_to<Key>, \equiv Why is this part of the interface? Why not Key::operator==()
  class Alloc = std::allocator< std::pair<const Key, T> >
  >
  class std::unordered_map;
```

What we need:

- A custom hash functor for case-insensitive strings
- A custom comparator functor, to compare strings ignoring character case

Exploring string hash tables <code walk-through>

```
template <class Type, class StringType = std::basic_string<Type>>
struct BasicStringHash
  using HashedType = StringType;
  size_t operator()(const HashedType & aStr) const
    std::hash<HashedType> hasher; ← we can use any hashing algorithm
    return hasher(aStr);
  }
  bool operator()(const HashedType & aStr1, const HashedType & aStr2) const
    return aStr1 < aStr2;</pre>
  struct KeyEquality
    bool operator()(const HashedType & aStr1, const HashedType & aStr2) const
      return aStr1 == aStr2;
  };
};
```

Exploring string hash tables <code walk-through>

typedef BasicStringHash<char> StringHash;

typedef BasicStringHash<wchar_t> StringHashW;

Eg.

std::unordered_map<wstring, TYPE, StringHashW, StringHashW::KeyEquality>

Exploring string hash tables <code walk-through>

```
template <class Type, class StringType = std::basic_string<Type>>
struct BasicStringHashI
```

```
using HashedType = StringType;
```

```
size t operator()(const HashedType & aStr) const
 // make a lower-case copy of the input string
 HashedType lowerStr(aStr);
 ToLower(const cast<Type *>(lowerStr.c str()));
 std::hash<HashedType> hasher;
 return hasher(lowerStr);
bool operator()(const HashedType & aStr1, const HashedType & aStr2) const
 return CompareI(aStr1, aStr2) < 0;</pre>
}
```

Case-insensitive hashes

//...

Exploring string hash tables <code walk-through>

```
template <class Type, class StringType = std::basic_string<Type>>
struct BasicStringHashI
{
    //...
    Case-insensitive hashes
    struct KeyEquality
    {
        bool operator()(const HashedType & aStr1, const HashedType & aStr2) const
        {
            return CompareI(aStr1, aStr2) == 0;
        };
    private:
```

```
static void ToLower(char * aStr) { ::CharLowerA(aStr); }
static void ToLower(wchar_t * aStr) { ::CharLowerW(aStr); }
```

```
static int CompareI(const string & aStr1, const string & aStr2)
{ return ::lstrcmpiA(aStr1.c_str(), aStr2.c_str()); }
```

```
static int CompareI(const wstring & aStr1, const wstring & aStr2)
{ return ::lstrcmpiW(aStr1.c_str(), aStr2.c_str()); }
```

Exploring string hash tables <code walk-through>

typedef BasicStringHashI<char> StringHashI;

typedef BasicStringHashI<wchar_t> StringHashWI;

Eg.

std::unordered_map<wstring, TYPE, StringHashWI, StringHashWI::KeyEquality>

Case-insensitive hashes



Show me the code!

Research Topic for You



Optimal file-path (case-insensitive)

hash functor for std::unordered_map<>

open4tech@caphyon.com

FilePath Hasher

Special type of case-insensitive hash: file-paths hash map.

(a hasher for the string representation of file paths, in a *case-insensitive* file system)

What we want:

std::unordered_map<FilePath, TYPE, FilePathHash, FilePathHash::KeyEquality>

Where FilePath encapsulates a std::wstring plus file specific methods.

What <u>issues</u> do we have with regular **StringHashWI** for file paths ?

std::unordered_map<FilePath, TYPE, StringHashWI, StringHashWI::KeyEquality>

FilePath Hasher

Requirements:

- The operations which have to be fast are *insertions* and *searches*
- Fast deletions would be desirable as well, but they are not a mandatory requirement
- A tradeoff between faster search time and slower insert time is accepted (if a faster search time can be achieved by slowing the insertion time a little bit)
- **FilePathHash** should be a drop-in replacement for **StringHashWI** (same API)
- **Benchmarks** for your FilePath hasher, having as baseline **StringHashWI**



So You Think You Can

Hashing Algorithms and Containers

inf.ucv.ro/~summer-school

Victor Ciura Principal Engineer



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