



UNIVERSITATEA  
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CRAIOVA

# Chasing Nodes

**Open4Tech: Graph Algorithms**

January 21, 2021



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**CAPHYON**



Q & A

Ask as we go along...

# Last Summer (2020)

## Open4Tech Summer School 2020

C++17/20 STL<Essentials>  
Code gold, not trash  
RESTful APIs

TikTok hand challenge recognition using Javascript  
Web Development Basics  
Processing web data with XML and XSLT



24 iunie - 10 iulie 2020  
<http://inf.ucv.ro/~summer-school/>

DEPARTAMENTUL DE INFORMATICA

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Open4Tech  
Summer School 2020  
24 iunie - 10 iulie  
ONLINE

## C++17/20 STL<Essentials>

Victor Ciura - Technical Lead



<http://inf.ucv.ro/~summer-school/>

[Slides: Open4Tech Summer School 2020 - C++ STL<Essentials>](#)

## STL Algorithms - Principles and Practice

*“Prefer algorithm calls to hand-written loops”*

*Scott Meyers, "Effective STL"*

# Why prefer to use (STL) algorithms?

👉 Goal: No Raw Loops }

*Sean Parent - C++ Seasoning, 2013*

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Whenever you want to write a **for/while** loop:

```
for(int i = 0; i < v.size(); ++i) { ... }
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# Why prefer to use (STL) algorithms?

👉 **Goal: No Raw Loops {}**

*Sean Parent - C++ Seasoning, 2013*

Whenever you want to write a **for/while** loop:

```
for(int i = 0; i < v.size(); ++i) { ... }
```

**Put the Mouse Down and  
Step Away from the Keyboard !**

*Burk Hufnagel*

# Why prefer to use (STL) algorithms?

## Correctness

Fewer opportunities to write bugs like:

- iterator invalidation
- copy/paste bugs
- iterator range bugs
- loop continuations or early loop breaks
- guaranteeing loop invariants
- issues with algorithm logic

# Why prefer to use (STL) algorithms?

## Code is a liability:

maintenance, people, knowledge, dependencies, sharing, etc.

**More** code => more bugs, more test units, more maintenance, more documentation

# Why prefer to use (STL) algorithms?

## Code Clarity

- Algorithm **names** say what they do
- Raw “for” loops don’t (without reading/understanding the whole body)
- We get to program at a higher level of **abstraction** by using well-known **verbs** (find, sort, remove, count, transform)
- A piece of code is **read** many more times than it’s **modified**
- **Maintenance** of a piece of code is greatly helped if all future programmers understand (with confidence) what that code does

# Why prefer to use (STL) algorithms?

## Simplicity

- Simpler code is more **readable** code
- Understandable and **expressive**
- Usually, **shorter** means simpler (*but not always*)
- **Unsurprising** code is more maintainable code
- **Idioms** are immediately recognized
- Code that moves complexity to abstractions (**libraries**) often has less bugs
- Compilers and libraries are often much better than you at **optimizing**
  - they're *guaranteed* to be better than someone who's not **measuring**

# Why prefer to use (STL) algorithms?

What's the difference?

## Performance / Efficiency

- Vendor implementations are highly **tuned** (most of the time)
- Avoid some unnecessary temporary copies (leverage **move** operations for objects)
- Function helpers and functors are **inlined** away (no abstraction penalty)
- Compiler optimizers can do a better job without worrying about pointer aliasing (auto-vectorization, auto-parallelization, loop unrolling, dependency checking, etc.)

# The difference between Efficiency and Performance

Efficiency	Performance
the amount of work you need to do	how fast you can do that work
governed by your algorithm	governed by your data structures

**i** **Efficiency** and **performance** are not necessarily dependent on one another.

# The difference between Efficiency and Performance

Why do we care ?

Because: *“Software is getting slower more rapidly than hardware becomes faster.”*

*“A Plea for Lean Software” - Niklaus Wirth*

lucid, systematic,  
and penetrating  
treatment of basic  
and dynamic data  
structures, sorting,  
recursive algorithms,  
language structures,  
and compiling

NIKLAUS WIRTH

**Algorithms +  
Data  
Structures =  
Programs**

PRENTICE-HALL  
SERIES IN  
AUTOMATIC  
COMPUTATION

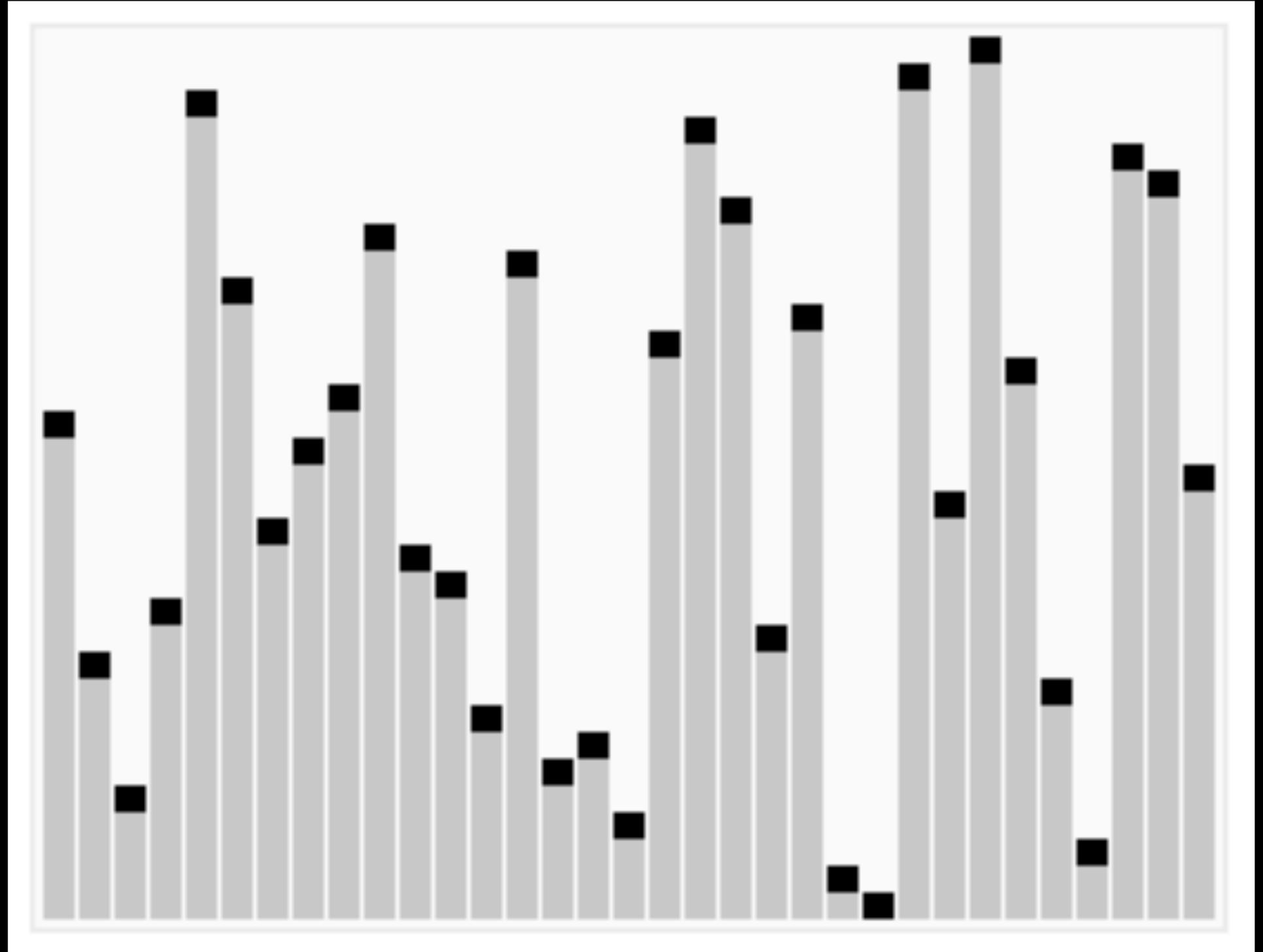
# The Big-O

Algorithm	Data structure	Time complexity:Best	Time complexity:Average	Time complexity:Worst	Space complexity:Worst
Quick sort	Array	$O(n \log(n))$	$O(n \log(n))$	$O(n^2)$	$O(n)$
Merge sort	Array	$O(n \log(n))$	$O(n \log(n))$	$O(n \log(n))$	$O(n)$
Heap sort	Array	$O(n \log(n))$	$O(n \log(n))$	$O(n \log(n))$	$O(1)$
Smooth sort	Array	$O(n)$	$O(n \log(n))$	$O(n \log(n))$	$O(1)$
Bubble sort	Array	$O(n)$	$O(n^2)$	$O(n^2)$	$O(1)$
Insertion sort	Array	$O(n)$	$O(n^2)$	$O(n^2)$	$O(1)$
Selection sort	Array	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(1)$
Bogo sort	Array	$O(n)$	$O(n n!)$	$O(\infty)$	$O(1)$

[wikipedia.org/wiki/Computational\\_complexity\\_theory](https://wikipedia.org/wiki/Computational_complexity_theory)

# The Big-O

Recognize the algorithm?

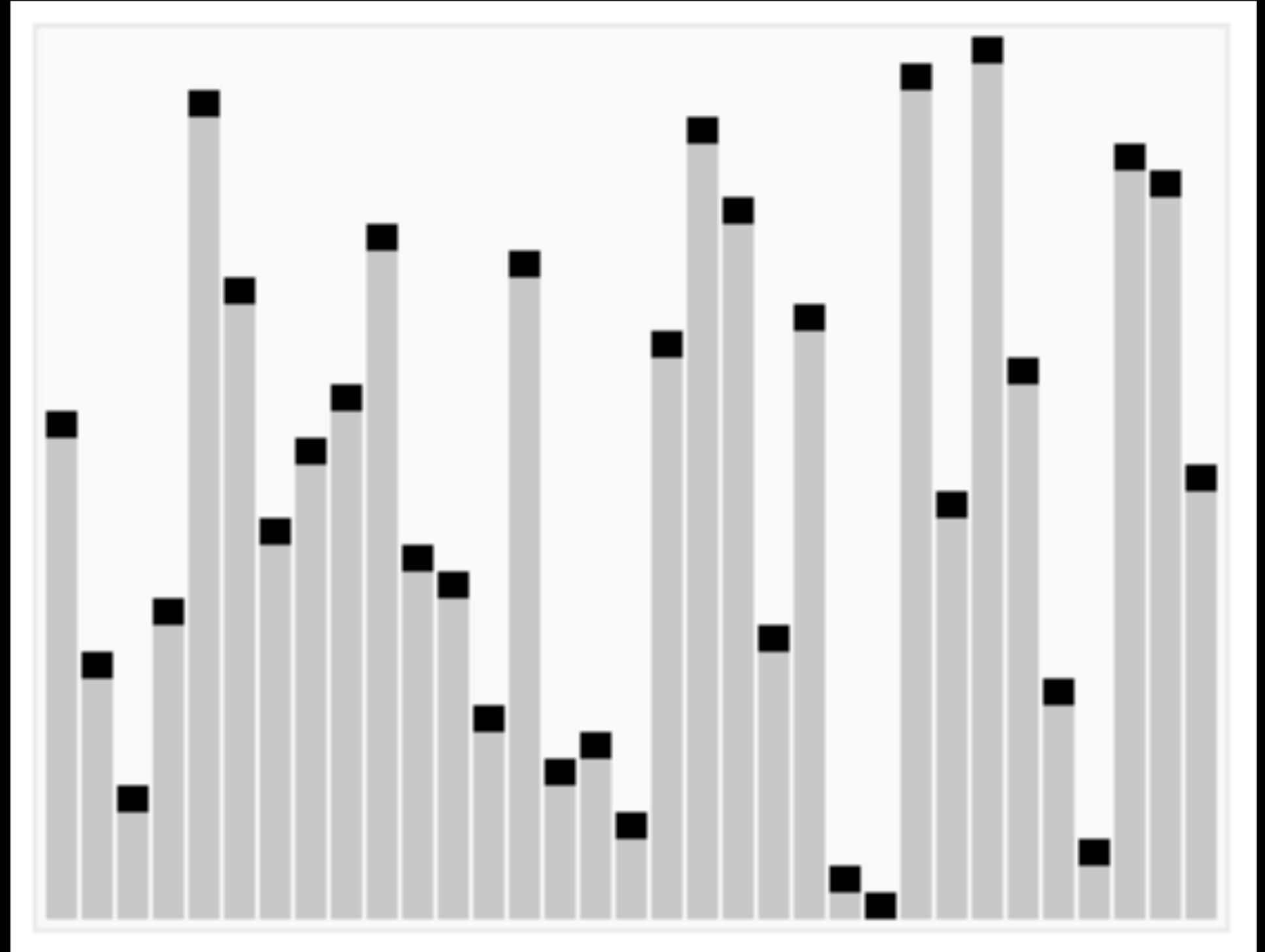




# The Big-O

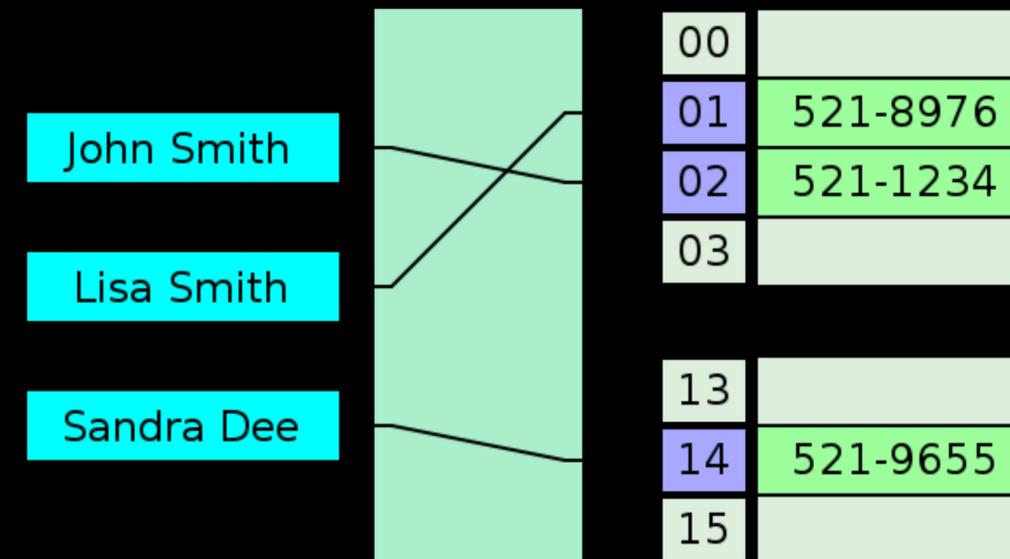
Recognize the algorithm?

**quicksort** algorithm  
has *average* case performance:  
 **$O(n \log n)$**



# What about Data Structures ?

Data structures along with the **operations** they provide, also have **complexity guarantees**



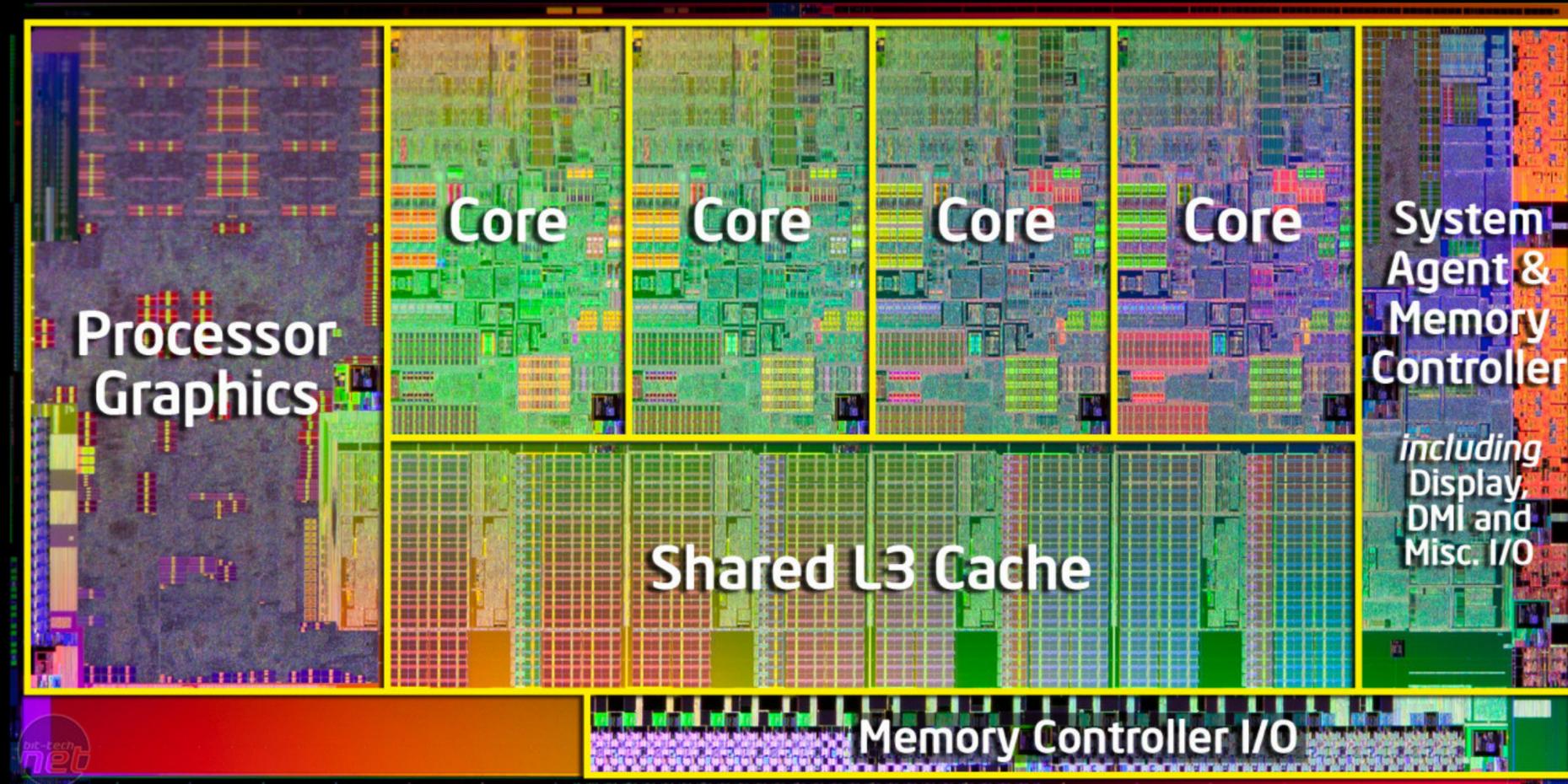
# STL Containers Big-O cheat-sheet

	A	B	C	D	E	F	G	H	I
1	<b>C++ STL</b>	<b>insert @end</b>	<b>insert @pos</b>	<b>erase @end</b>	<b>erase @pos</b>	<b>find</b>	<b>sort</b>	<b>iterator</b>	<b>comment</b>
2	<b>vector</b>	0(1)	0(dist(pos,end))	0(1)	0(dist(pos,end))	0(n)	0(n*log(n))	RandomAccess	array
3	<b>deque</b>	@begin/@end 0(1)	0(dist(pos,begin/end))	@begin/@end 0(1)	0(dist(pos,begin/end))	0(n)	0(n*log(n))	RandomAccess	
4	<b>list</b>	0(1)	0(1)	0(1)	@pos 0(1); @key 0(n)	0(n)	0(n*log(n))	Bidirectional	doubly linked
5	<b>stack</b>	0(1) push()	-	0(1) pop()	-	0(n)	-	same as container	adaptor<deque, list, vector>
6	<b>queue</b>	0(1) push()	-	0(1) pop() @begin	-	0(n)	-	same as container	adaptor<deque, list>
7	<b>set/map</b>	-	0(log(n))	-	@pos 0(1); @key 0(log(n)+count(key))	0(log(n))	sorted	Bidirectional	red-black tree (balanced BST)
8	<b>unordered_set/ unordered_map</b>	-	avg 0(1); worst 0(n)	-	@pos avg 0(1) worst 0(n); @key 0(count(key))	avg 0(1); worst 0(n)	-	Forward	hash_set/hash_map
9	<b>priority_queue</b>	push() 0(log(n))	-	pop() 0(log(n))	-	top() 0(1)	-	RandomAccess	adaptor<vector, deque> => constant time extraction of the largest (default) element, at the expense of logarithmic insertion
10	<b>make_heap(range)</b>	push_heap() 0(2*log(n))	-	pop_heap() 0(2*log(n))	-	max is first	0(n*log(n))	RandomAccess	constructs a max heap in the range

# What about Performance ? 🚀

How **fast** can the CPU execute **each step** from the algorithms.

This is mostly determined by the native (CPU) **data types** used and *your choice* of **data structures**.



## Strategy

- **Identification:** *profile* the application and identify the worst performing parts
- **Comprehension:** understand what the code is trying to achieve and why it is slow
- **Iteration:** change the code based on step 2 and then re-profile; repeat until fast enough

Don't trust your instinct !  
Always **benchmark** the code changes.

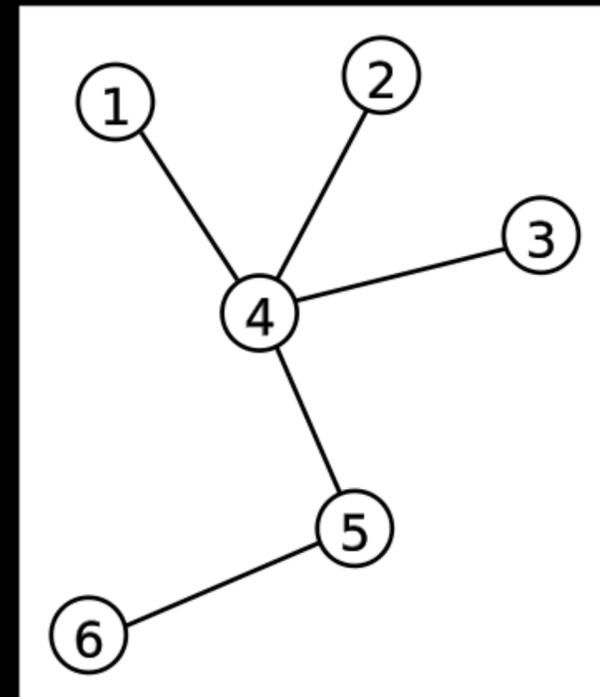
Very often, code becomes a **bottleneck** for one of four reasons:

- It's being called too often
- It's a bad choice of algorithm:  $O(n^2)$  vs  $O(n)$ , for example
- It's doing unnecessary work or it is doing necessary work too frequently
- The data is bad: either too much data or the layout and access patterns are bad

## Today, let's focus on **data structures**

Because this is part of a course on graph algorithms,

let's focus specifically on *node-based* data structures: **graphs & trees**.

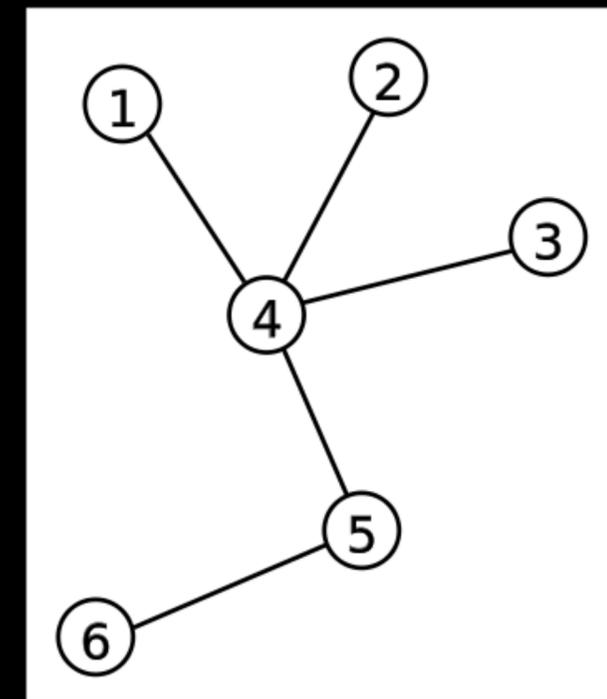


# Focus > Narrowing

In graph theory,

**tree** is an **undirected graph** in which any two vertices are connected by exactly one path, or equivalently a connected acyclic undirected graph.

**forest** is an **undirected graph** in which any two vertices are connected by at most one path, or equivalently an acyclic undirected graph, or equivalently a disjoint union of trees.



## Tree data structures

Abstract data type that simulates a hierarchical tree structure, with a **root** value and **subtrees of children** with a parent node, represented as a set of **linked nodes**.

# Trees

You probably already know a lot about trees, of different types, each with individual **specific properties** and **use cases** in computer science.

# Trees

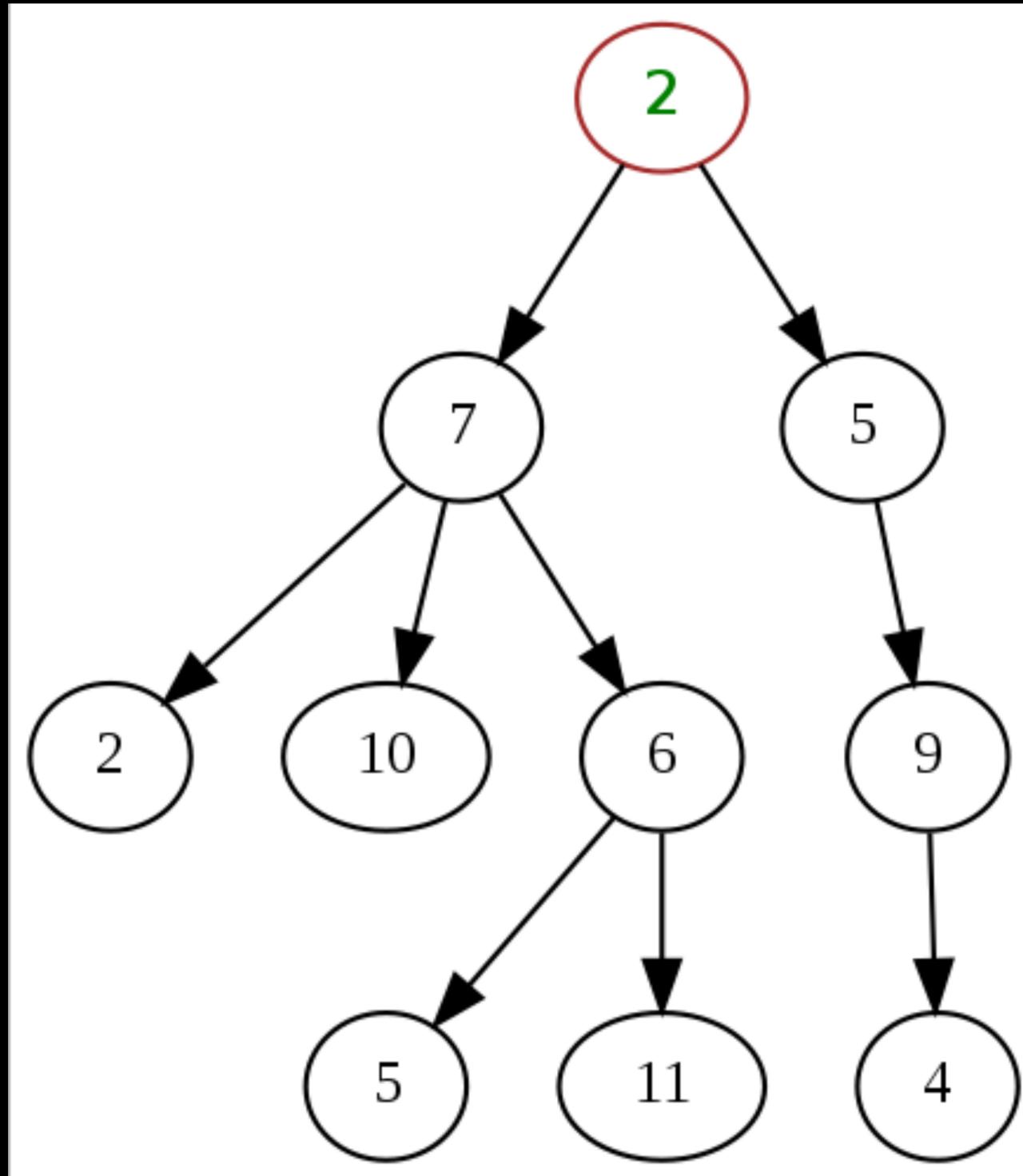
You probably already know a lot about trees, of different types, each with individual **specific properties** and **use cases** in computer science.

But, I bet you don't know they look like this:

# Trees



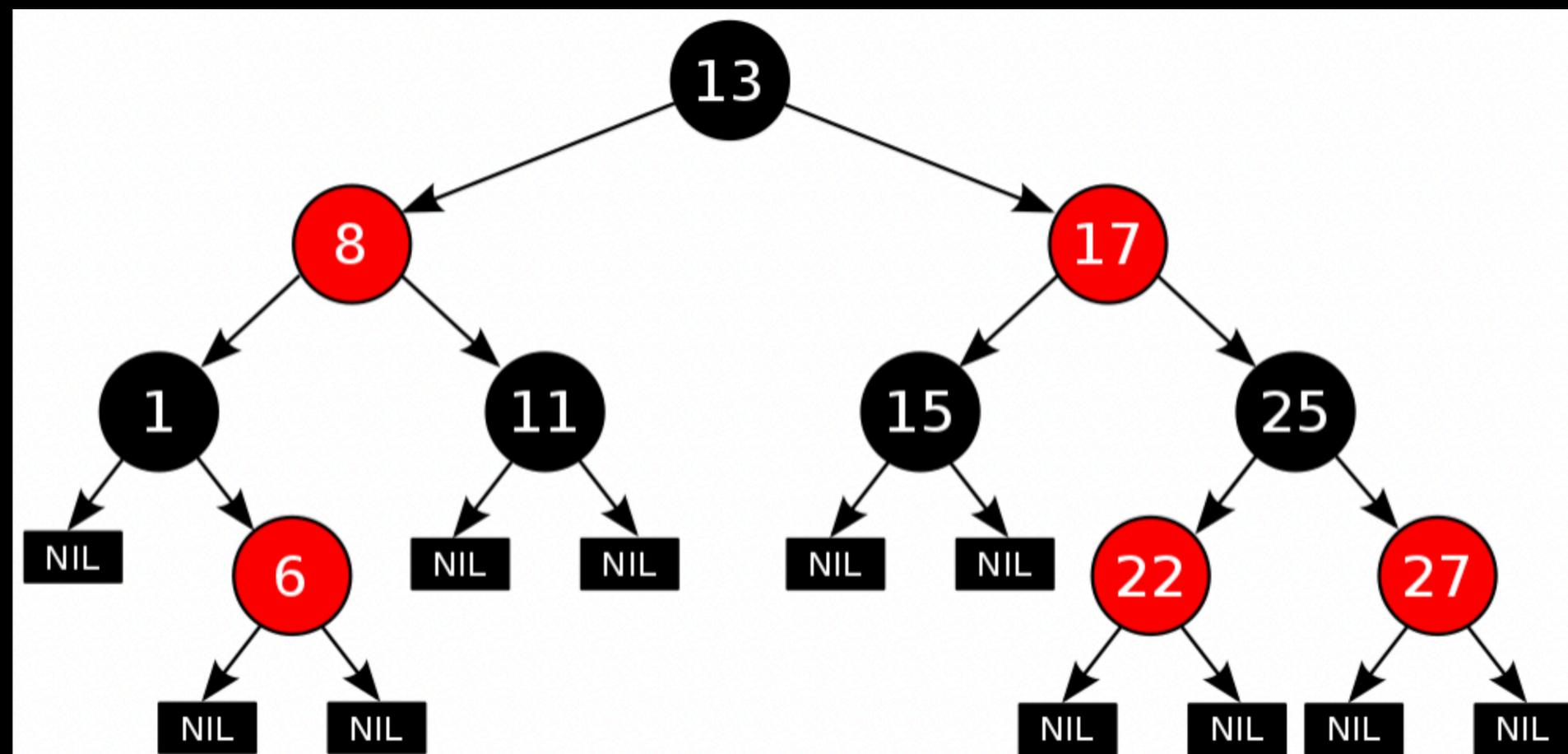
# Trees



## Red-black tree

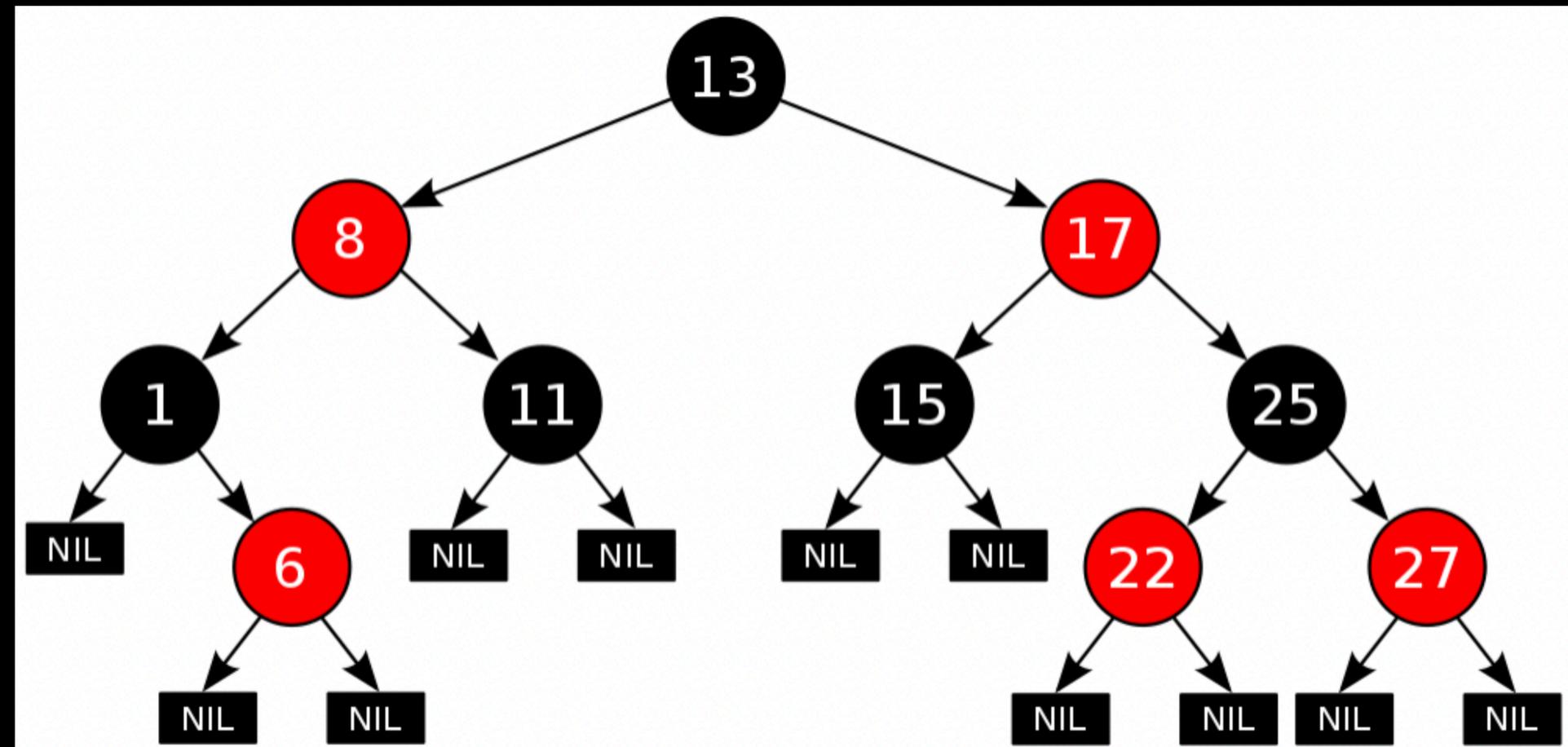
Self-balancing **binary search tree**

Each node stores an extra bit representing **color** (red/black), used to ensure that the tree remains *approximately* balanced during insertions and deletions.



# Red-black tree

	Average/Worst
space	$O(n)$
lookup (search)	$O(\log n)$
insert	$O(\log n)$
delete	$O(\log n)$



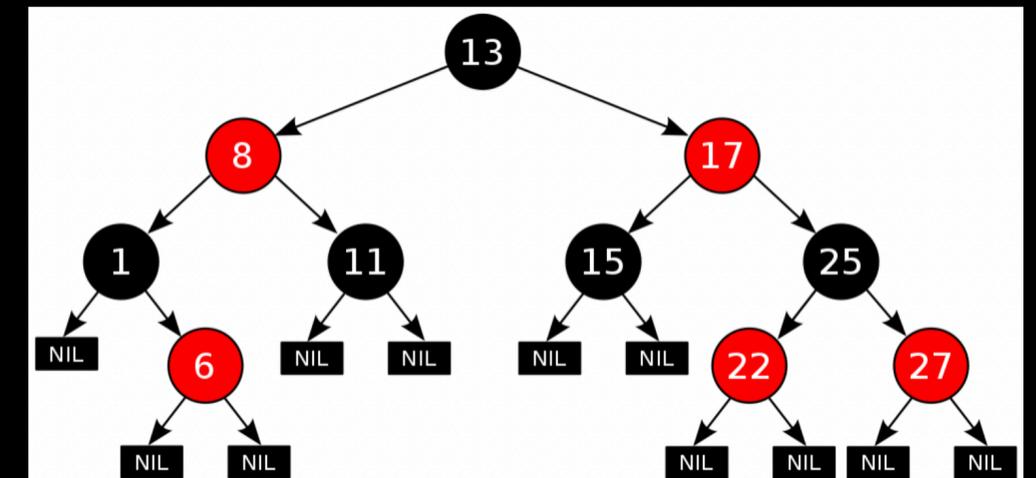
As opposed to other BSTs, the **re-balancing is not perfect**, but guarantees searching in  $O(\log n)$  time

# Red-black tree

**Why** am I narrowing to this special kind of binary search tree?

Because **Alex Stepanov** picked this kind of tree as the reference implementation for the API he designed for C++ STL associative data structures:

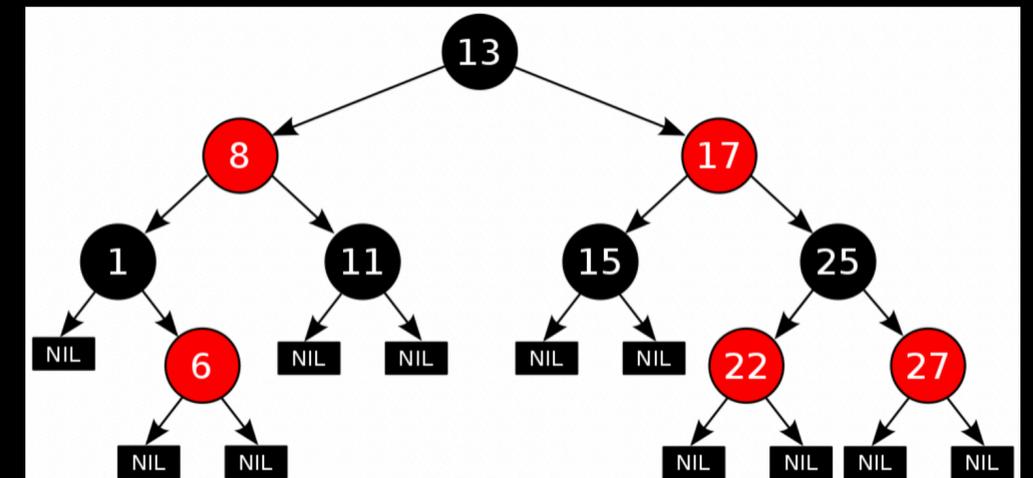
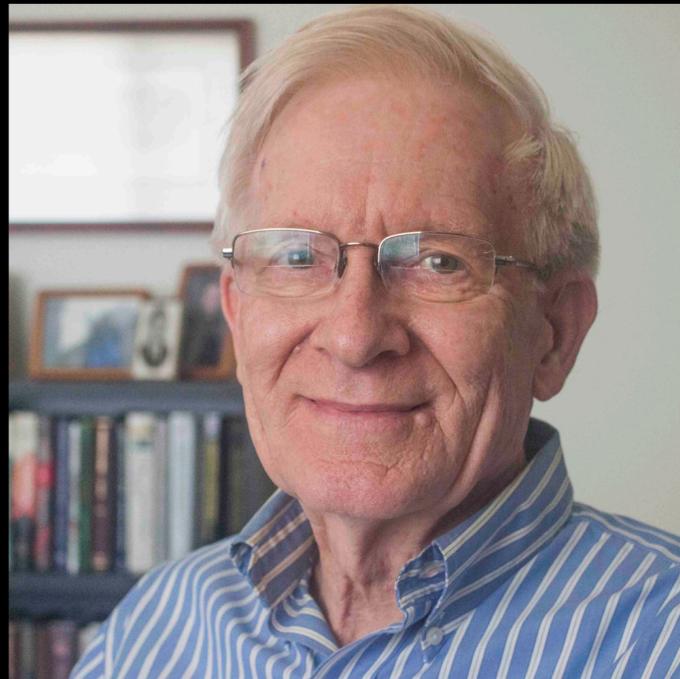
**`std::map`** & **`std::set`**



# Red-black tree

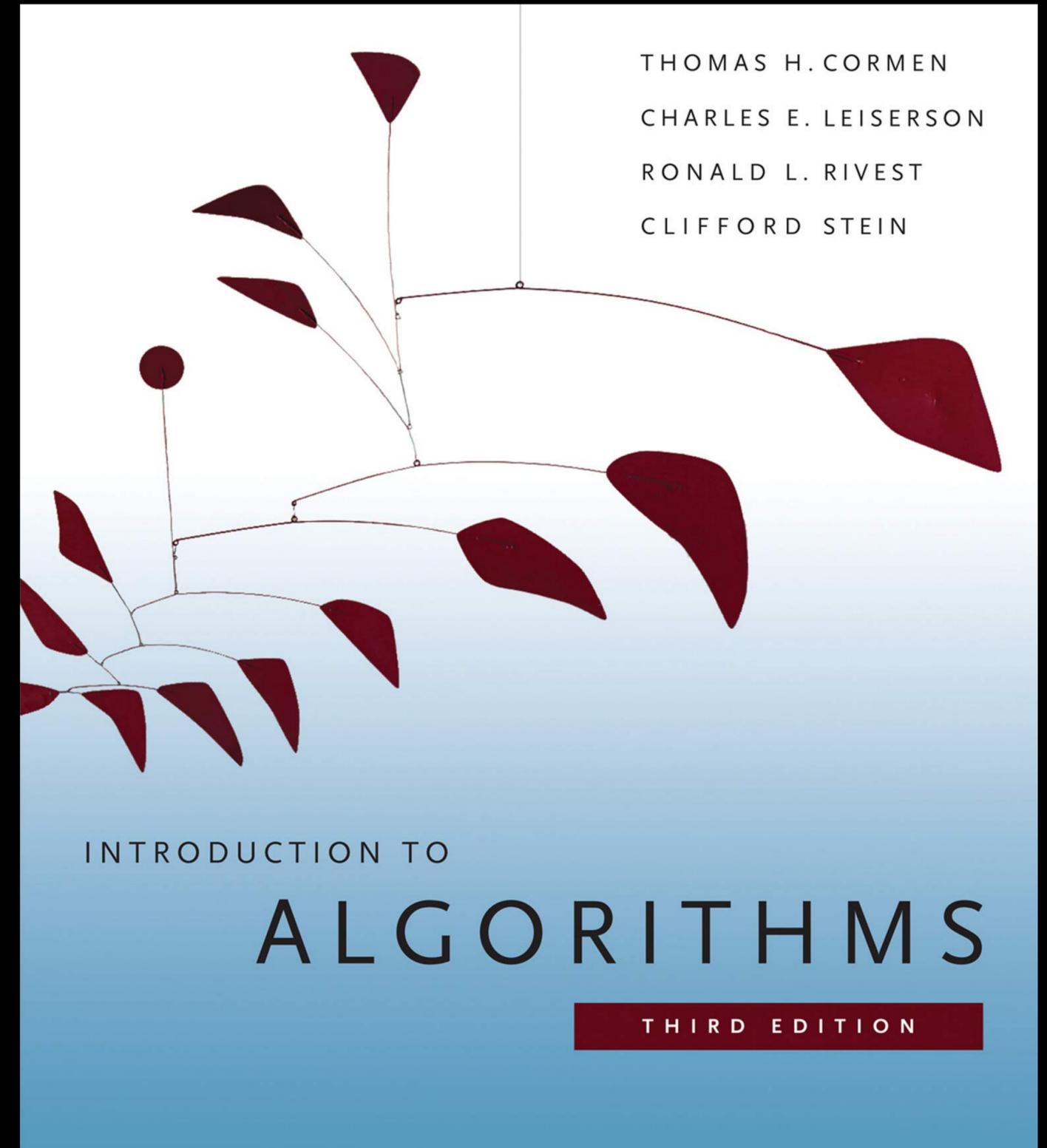
**Why** am I narrowing to this special kind of binary search tree?

**David Musser** coded the best C++ *implementation* for the API Stepanov designed:  
**`std::map`** & **`std::set`**



# The Book

If you want to dig deep,  
I highly recommend this classic:



# Red-black tree

Red-black trees are very **advanced** data structures, that are beautifully wrapped in a very easy to use API:

**std::map** & **std::set**



# Red-black tree

Red-black trees are very **advanced** data structures, that are beautifully wrapped in a very easy to use API:

`std::map` & `std::set`

... and this is where things get interesting 🤪

Let's see!



# Code dive

We'll explore together these properties, by building a **search engine index** in C++

... Really 😄

Let's see what we want to build.

# Search engine index

## Google Autocomplete

As you type in the browser search box, you can find information quickly by seeing **search predictions** that might be similar to the search terms you're typing.

The suggestions that Google offers all come from how people actually search.

For example, type in the word “cruise” and you get suggestions like:

Keyword: **cruise**

Suggested searches for: **"cruise"**

- > cruise line
- > cruise ship
- > carnival cruise
- > caribbean cruise
- > princess cruise
- > disney cruise
- > celebrity cruise
- > norwegian cruise
- > alaska cruise
- > ship cruise

# Search engine index

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These are all real searches that have been done by other people.

**Popularity** is a factor in what Google shows.

If lots of people who start typing in "**cruise**" then go on to type "**line**" that can help make "**cruise line**" appear as a suggestion in the future.

# The task

We have a keywords “database” in the form of a large text file ( [keywords.db](#) ) containing search terms (phrases) used by people in the past.  
(consider this an *active search cache*)

Here is a *small fragment* from this text file:

```
----- keywords.db -----  
philips lcd 15  
15 lcd cheap monitor  
cheap 15 lcd monitor  
dell e153fp 15 lcd midnight grey 36  
lcd tv 15  
samsung lcd 15  
sony 15 lcd monitor  
15 dvd lcd tv  
15 inch lcd plasma monitors  
...
```

# Assumptions

You may assume the following simplifying **preconditions**:

- the text file contains only **ASCII** alphanumeric characters (English words)
- keywords are separated by **space** or **CR/LF**
- keywords database file is to be considered an **immutable** (read-only) snapshot of the current query cache
- each line in the file represents a **search phrase** used in the past
- consider the whole “database” as a **continuous** chain of keywords, separated by whitespace
- a **keyword** is a sequence of non-whitespace characters (words)

# Search phrase

For simplicity, we shall define a **search phrase** as a pair of just **two** consecutive **keywords** in the query database.

E.g.

"cruise line"

"dell e153fp"

"cruise ship"

"samsung lcd"

"norwegian cruise"

"lcd cheap"

"sony 15"

"cheap monitor"

# First task

First, we have to **load** and **rank** the keywords database.

That means **ordering** all search phrases according with their **frequency** in the cache (database).

We should be able to print the **Top 1000** search phrases with their respective **ranks** (occurrence frequency).

# First task

E.g.

Top 10 search phrases from **keywords.db** with their respective **# ranks**

```
real estate # 43298
for sale # 38022
new york # 27302
how to # 25068
web site # 21073
las vegas # 19039
cell phone # 17657
of the # 15012
credit card # 14278
web hosting # 11037
```

# Second task

Our second task is to implement our own **auto-suggestion engine** for 10 related searches, based on top search phrases containing the **input** keyword.

See previous *example* with suggested searches for keyword: "cruise".

This operation should be **super-fast**.

\* This **interactive mode** should be active only when the program receives a **/search** command-line switch.

# The Code

We're going to see **2** completely **different implementations** for this program.

We're going to analyze the PROs & CONs of each and see some **hints** for a potential **3rd** implementation => your **homework** assignment.



[open4tech@caphyon.com](mailto:open4tech@caphyon.com)

# Solution 1

## Data Structures

Data structures used by the algorithm are designed to store the **minimal** amount of information in **memory** (no redundancy, no keyword copies).

Data structures leverage STL container **iterators** that are **stable** (valid) under the used algorithm operations.

We use **node-based** data structures (red-black trees): **`std::set`** & **`std::map`**

# Solution 1

## The Algorithm

Loading the keyword database into our data structures (counting search phrase occurrences).

=> filling a `std::map` from each phrase combination to its `frequency`

=> using `std::set` & `std::map` iterators everywhere, to avoid copying strings (keywords)

=> keywords are stored & referenced from a single location in an `std::set` (unique)

=> `ranking` is done automatically by means of a custom `std::set` comparator `predicate`

## DEMO TIME

Let's dive into the code...

# Solution 1 - analysis

**PROs**

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## PROs

- is a very good **showcase** for STL usage (serves its **didactical** purpose)

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- is **idiomatic** STL usage

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- offers good **memory** working set scaling for **long search phrases**

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- for simplicity, our implementation uses **case-sensitive compare** for keywords

# Solution 2

## Data Structures

Are not designed to store the minimal amount of information in memory, having considerable **redundancy** in storing the keywords (allows for storing **duplicate** instances of keywords).

We use an STL **unordered\_map** container to store all search phrases and their occurrences.

We store each keyword pair as a *concatenated* string "**keyword1 keyword2**" (map-first) with its corresponding **counter** (map-second).

This is where our data redundancy stems from (**duplicated keywords** from search pairs).

# Solution 2

## Data Structures

We chose this advanced data structure for our algorithm, because it is a **hash map**. We leverage this fact for its **speed** in storing a new search phrase and finding an existing tuple to increment its frequency (**in constant time**).

Usage of the **CompareKeywordTupleCount** custom binary **predicate** is **optional**, because it is not mandatory to perform a **stable sort** (*lexicographic*) with regards to search phrases (keyword pairs) that have the same rank/frequency.

# Solution 2

## The Algorithm

Loading the keyword database into our data structures (counting search phrase occurrences).

=> filling a `std::unordered_map` from each phrase combination to its frequency

```
"keyword1 keyword2" # 24
```

=> ranking keyword database using an auxiliary `std::vector` and applying `std::sort()` algorithm with a *custom predicate* (lexicographic stable sort, optional)

## DEMO TIME

Let's dive into the code...

# Solution 2 - analysis

**PROs**

# Solution 2 - analysis

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- is **idiomatic** STL usage
- is **type-safe** and **memory safe**
- offers **good performance** characteristics for large data sets
- it's **very fast** (due to hash-based **lookup**)
- although it **duplicates** data, its **memory usage is lower** than [Solution 1], because we have short keywords in our database and [Solution 1] has a lot of memory waste due to tree node 64-bit pointers

# Solution 2 - analysis

**CONS**

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- it stores **duplicated** keywords (cannot help but feel uncomfortable about this ?!)

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# Solution 2 - analysis

## CONs

- it stores **duplicated** keywords (cannot help but feel uncomfortable about this ?!)
- offers **poor memory working set scaling** for long search phrases  
(due to data duplication)
- for simplicity, our implementation uses **case-sensitive** compare for keywords
- it uses the notoriously **slow I/O streams** for data input

# Solution 3 - Hints

Alternative solutions and further improvements:

- We could use a **memory mapped file** to map the keyword database directly into process memory, so that we could avoid using I/O streams and string parsing, processing
- We could perform a **partial\_sort** of the keyword tuples (just **Top N** search phrases) and perform our lookup for suggestions in that pool

# Solution 3 - Hints

Alternative solutions and further improvements:

- We could use a much more **cache-friendly** data structure, like an **std::vector** to store the tuple counts more compactly (array).
  - we would sort the array
  - count adjacent equal pairs
  - store counts and tuples in another array that we (partially) sort
  - read out the range desired

# Solution 3 - Hints

Alternative solutions and further improvements:

- Because we are dealing strictly with **English** words, we could cut off (**truncate**) keywords at **8 bytes each** and store them in a **uint64\_t** integer.

⚠ This is not functionally equivalent, but good enough because **most keywords** in the database are smaller than 8 characters.

Using **integers** instead of strings would be a huge performance boost when performing comparisons and would also be much more space efficient.

# Solution 3 - Hints

If you want to try using these hints to build an even better solution, I want to see it 😊



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**HAVE FUN !**



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# Chasing Nodes

**Open4Tech: Graph Algorithms**

January 21, 2021



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